

Optimal performance of heterogeneous networks based on the bit rate

Déthié Dione¹ · Salimata Gueye Diagne¹ · Bakary Koné¹ · Youssou Gningue²

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Abstract Networks LTE(4G) and Wi-Fi complementarity establish a heterogeneous system of wireless and mobile networks. We study and analyze the optimal performances of this heterogeneous system based on the bit rate, the blocking probability and user connection loss. Random Waypoint is the user mobility model. User's provided mobile terminal equipped with multiple accesses interfaces. We have developed a Markov chain to estimate the performances obtained from the heterogeneous networks system, which allowed us to propose an average bit rate value in a sub-zone of this system then the average blocking probability user connection in this zone. We have also proposed a sensitivity factor of maximal decrease of these selection network parameters. This factor informs about the heterogeneous networks congestion and dis-congestion system.

Keywords LTE · Wi-Fi · Random waypoint · Handover · Markov chain · BLER · Bit rate · Blocking probability

Introduction

The integration of wireless and mobile networks such as longterm evolution (LTE) and Wi-Fi is nowadays a necessity for the satisfaction of user request which is stronger and stronger. The global services and user mobility makes this task difficult. However these networks are provided with characteristics to support services and user mobility.

The Random waypoint (RWP) is model chosen for users mobility in this heterogeneous networks system. The RWP model corresponds to the ideal behavior of a user in an urban area because according to the model of mobility RWP, every user chooses randomly a place of destination and goes to this one in a constant velocity. The user movement starting up from a point to a destination is named "one movement epoch". The users velocity in every epoch is random variable and is chosen from a uniform distribution of the velocity $[0, V_{max}]$, where V_{max} is the eligible maximal velocity for user. In the RWP, the user can wait for a period of time called "time of reflection" before his departure for another point. The browsed path of a user is independent from its previous path and from other users path. Thus, at the end of every time of movement, the user stops a duration of the time and then chooses another destination place and, possibly, new velocity, and moves towards this destination in a constant velocity, and so on. RWP is one of the mobility models widely used in the performances analysis of the wireless and mobile networks. It represents well the individual movements which include the stop, the starting up and the other actions bound to the individual movements in cities.

So new methods for saving, transmission and sharing of bandwidth are imperative. Among these we can mention the selection technique of best network based on the bit rate which the selection parameters of which are the blocking probability and connections losses.

We noted in the literature the most used selection strategies of network. In their works [10], we were able to analyze the signal power received (RSNS) and then the available bandwidth (TBNS) of a heterogeneous networks system. They emphasized the parameters of this system such as the blocking probability and connections losses but they did not take into account the interference in the selection techniques

✉ Déthié Dione
dethiedione79@gmail.com

¹ Department of Mathematic and Computer Science,
University of Cheikh Anta Diop, Dakar, Senegal

² University of Laurentia, Sudbury, Canada

that they developed which made less successful results obtained from the blocking probability and connections losses. Besides the authors [1, 2, 12, 13] considered the interference in the selection strategies they adopted which is the one based on the SINR, which allowed them to decrease the probability connections losses during a vertical handover. But through their studies, they did not approach the parameters connections blocking. On the other hand, these are analyzed by [7] who in their works obtained results more satisfactory than those who are preceded them. Besides, [11] took into account the users mobility (terminal-controlled mobility management) and other aspects such as the cost, the battery life cycle and the handover frequency.

However, through their studies, none of these authors took into account the constraint related to the block error rate (BLER) in a sub-cell given by the cluster. The contribution in our works bases on our model users connected and disconnected. When he is connected, the system quits a well-defined state and moves to an other state before returning to this first one in a given time interval. Besides, unlike the other works, we took the bit rate as an important selection technique of the best network which consists in choosing the biggest bit rate value.

The paper is organized as follows: “[Model of heterogeneous networks system](#)” section introduces the model of the studied heterogeneous networks where all the parameters of the system are defined. The selection method algorithm based on the bit rate is established at the level of “[Selection method](#)” section. In “[User mobility model](#)” section, we have developed the users mobility model, it is Random Waypoint (RWP) which we consider more adequate to the individual users movements. The average access demand rate for a service is given in “[Average new access demand rate \$\lambda_{C_i}^{C\(k\)}\$ for a service](#)” section. The average rate demand of vertical and horizontal handover is calculated in “[Average demand rate of handovers](#)” section. In “[Modeling approach based on a Markov chain](#)” section, we have used a Markov chain to analyze the number busy bandwidth units. The system studied performances such as the bit rate and the blocking probability and connection losses are estimated in “[Evaluations of optimal performances](#)” section. The results obtained are simulated in “[Numerical tests](#)” section. We finished our study in “[Conclusion](#)” section by conclusion and future work.

Model of heterogeneous networks system

The model of heterogeneous networks system which we study is represented by Fig. 1.

Indeed, we have an hexagonal area of service C_1 covered entirely by the mobile network LTE (4G). In this

service area are present several homogeneous circular sub-cells $(C_j)_{2 \leq j \leq m}$ of radius r_i among which each is also covered by a Wi-Fi wireless. So both mobile network LTE and Wi-Fi wireless overlap in cells C_i and the Wi-Fi wireless are separated between them. We denote by C_0 the part of the service area not covered by a Wi-Fi wireless. Where from we have:

$$C_0 = C_1 - \bigcup_{j=2}^m C_j \quad (1)$$

As the users are equipped with devices of multiple accesses, they have the possibility to connect or disconnect from a network in the cells where networks overlap by choosing automatically the network which has the best bit rate. In our study, we suppose that the LTE network supplies two types of services: those Multicasts or Unicasts whose numbers of units of bandwidth are respectively B_1^{mc} et B_1^{uc} . Besides, the number of units of bandwidth for every Wi-Fi wireless is B_i .

In the service area C_1 , we suppose to have Q interferences sources distributed following a normal random distribution: $Q = \{I(q), q = 1, \dots, Q\}$.

The selection technique is based on the bit rate then these interferences play an important role at the level of

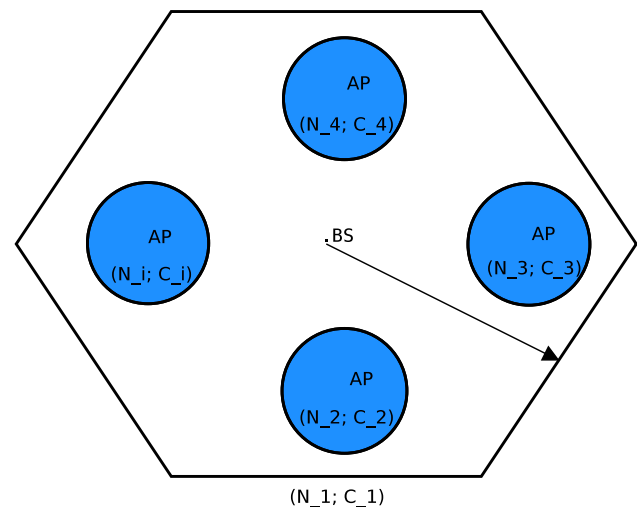


Fig. 1 Model cluster service zone

Table 1 LTE network parameters

Parameters	LTE networks N_1	
Interferences	Q sources	
Covered zone	C_1	
Service quality	Low bandwidth	
Types of services	Multicast	Unicast
Units of bandwidth	B_1^{mc}	B_1^{uc}

The arrival rate $\tau_a(C_i)$ of a user in a sub-cell C_i

By means Random WayPoint model and results obtained by [4], we have proved that the average arrived rate of a user in a sub-cell C_i of radius r_i and situated at a distance d_i of the cluster center is established by equation:

$$\tau_a(Z_i) = \frac{2}{C_v} \int_0^\pi \int_0^\pi r_i \cdot h(x) \sin \Phi d\Phi dz \quad (3)$$

Probability $\mathbb{P}(C_i)$ of finding users in a sub-cell C_i

The probability of finding users in a cell C_i of radius r_i situated at a distance d_i of the service area center depends directly on the users mobility. Thus we are based on the Random WayPoint and works results of the authors [3] to show this probability by equation:

$$\mathbb{P}(C_i) = \int_0^R \int_0^\pi \int_0^{2\pi} r_i \cdot h(x) d\Phi dz dr_i \quad (4)$$

Average new access demand rate $\lambda_{C_i}^{C(k)}$ for a service

Let us denote by $\lambda_{C_1}^{C(k)}$ the average demand access rate for a service k in a cell C_1 . So the average demand access rate for a service k , $\lambda_{C_i}^{C(k)}$, being in one sub-cell C_i is defined by formula:

$$\lambda_{C_i}^{C(k)} = \mathbb{P}(C_i) \cdot \lambda_{C_1}^{C(k)} \quad (5)$$

Average demand rate of handovers

A connected user and then in mobility makes indirectly handover in the heterogeneous networks system. This handover is horizontal or vertical.

Horizontal

Let $u_{C_0}^k$ the average user number having access to the service k in the cell C_0 . The average demands rate of horizontal handover $\tau_{C_0}^{H(k)}$ to the network N_1 for a service k is given by relation 6 referring to the author (Jabban et al. [6]):

$$\tau_{C_1}^{H(k)} = u_{C_0}^k \cdot \eta_{C_0}^{C_1} \quad (6)$$

where $\eta_{C_0}^{C_1}$ is the users exit flow of the zone C_0 outside cell C_1 and is defined by:

$$\eta_{C_0}^{C_1} = \frac{\mathbb{P}(C_0)}{\Delta_{C(k)}}$$

with $\Delta_{C(k)}$ the average residence time of users in a zone C_i .

Vertical

We denote by $u_{C_0}^k$ the average number of mobile users who have accessed to the service k in the cell C_0 . The average demands rate of vertical handover $\tau_{C_0}^{V(k)}$ to the network N_i of users who have accessed to the service k in the zone C_0 and moving towards the cell C_i without finishing their connections is defined by Eq. 7 refer to the author (Jabban et al. [6]):

$$\tau_{C_i}^{V(k)} = u_{C_0}^k \cdot \eta_{C_0}^{C_i} \quad (7)$$

where $\eta_{C_0}^{C_i}$ is the users exit flow of the cell C_0 towards cell C_i and is given by:

$$\eta_{C_0}^{C_i} = \frac{\mathbb{P}(C_0)}{\Delta_{C(k)}}$$

with $\Delta_{C(k)}$ the average residence time of the users in a zone C_i .

Modeling approach based on a Markov chain

We have developed a Markov chain for modeling the dynamic fluctuations and define all the stages and states of the heterogeneous networks system.

Supposing that the set cells and services which are present in the service area, are respectively denoted M and S so the Markov chain size is given by : $s.(2m + 1)$ with $|M| = m$ and $|S| = s$.

Different stages and states of the system

When we take the system at the given moment then we define it as being a stage of dynamic change.

Besides, user connections and disconnections from a network give the various states of the system which define the states space by (see Table 3):

$$\mathcal{E} = \{(b_{1,1}^k; b_{1,2}^k; \dots; b_{1,i}^k; \dots; b_{1,m}^k; b_2^k; \dots; b_i^k; \dots; b_m^k)\}$$

$$s.q./ \left\{ \begin{array}{l} \sum_{j=1}^m \sum_{k=1}^s (b_{ij}^k) \leq B_1^{uc} \\ \sum_{k=1}^s (b_i^k) \leq B_i \end{array} \right.$$

States are differentiated by the possible variation of the units of busy bandwidth in a given zone. For example when the system is in a given state, it changes state if a user connects recently or to make a handover either disconnects by freeing units of busy bandwidth.



Table 3 Definition of system parameters

Parameters	Definitions
$b_{1,1}^k$	Number of units of busy bandwidth of LTE network to the service k in the cell C_0
$b_{1,i}^k$	Number of units of busy bandwidth of LTE network to the service k in the cell C_i
b_i^k	Number of units of busy bandwidth of LTE network to the service k in the cell C_i
N_{PRB}^k	Number of resources blocks asked to supply a service k by the LTE network in the cell C_i
$\sum_{k=1}^s b_{1,1}^k = b_{1,1}$	Number of units of busy bandwidth of LTE network in the cell C_0
$\sum_{k=1}^s b_{1,i}^k = b_{1,i}$	Number of units of busy bandwidth of LTE network in the cell C_i
$\sum_{k=1}^s b_i^k = b_i$	Number of units of busy bandwidth of Wifi wireless in the cell C_i

$$b_{11} \longrightarrow b_{11} + N_{PRB}^k$$

Where b_{11} is the number of units of busy bandwidth in the zone C_1 by the network N_1 and the number of busy blocks resources by a user for a service k in the zone C_1 to the network N_1 .

- Stage 0: $E_0 = (b_{1,1}^k; \dots; b_{1,i}^k; \dots; b_{1,m}^k; b_2^k; \dots; b_i^k; \dots; b_m^k)$
- Stage 1: $E_1 = (b_{1,1}^k + \mu N_{prb}^k; b_{1,2}^k; \dots; b_{1,m}^k; b_2^k; \dots; b_m^k)$

$$\mu = \begin{cases} 1 & \text{if a user } u_0 \text{ connects to the network LTE in the cell } C_0 & \text{State}(E_{1,1}) \\ -1 & \text{if a user } u_0 \text{ disconnects of the LTE from the cell } C_0 & \text{State}(E_{1,2}) \\ 0 & \text{Otherwise} & \text{State}(E_{1,3}) \end{cases}$$

- Stage 2: $E_2 = (b_{1,1}^k; \dots; b_{1,i}^k + \mu N_{prb}^k; \dots; b_{1,m}^k; b_2^k; \dots; b_m^k)$

$$\mu = \begin{cases} 1 & \text{if a user } u_0 \text{ connects to the network LTE in the cell } C_i & \text{State}(E_{2,1}) \\ -1 & \text{if a user } u_0 \text{ disconnects of the network LTE from the cell } C_i & \text{State}(E_{2,2}) \\ 0 & \text{Otherwise} & \text{State}(E_{2,3}) \end{cases}$$

- Stage 3: $E_3 = (b_{1,1}^k; \dots; b_{1,m}^k; b_2^k; \dots; b_i^k + \mu; \dots; b_m^k)$

$$\mu = \begin{cases} 1 & \text{if a user } u_0 \text{ connects to the wireless Wi-Fi in the cell } C_i & \text{State}(E_{3,1}) \\ -1 & \text{if a user } u_0 \text{ disconnects of the wireless Wi-Fi from the cell } C_i & \text{State}(E_{3,1}) \\ 0 & \text{Otherwise} & \text{State}(E_{3,1}) \end{cases}$$



- Stage 4: $E_4 = (b_{1,1}^k + \mu N_{PRB}^k; \dots; b_{1,i}^k + \mu' N_{PRB}^k; \dots; b_{1,m}^k; b_2^k; \dots; b_m^k)$

$$(\mu, \mu') = \begin{cases} (1, -1) & \text{if a user } u_0 \text{ connects to the network LTE in the cell } C_0 \\ & \text{by disconnecting of network LTE from the cell } C_i & \text{State}(E_{4,1}) \\ (-1, 1) & \text{if a user } u_0 \text{ disconnects of the network LTE from the cell } C_0 \\ & \text{by connecting to the network LTE in the cell } C_i & \text{State}(E_{4,2}) \\ (0, 0) & \text{Otherwise} & \text{State}(E_{4,3}) \end{cases}$$

- Stage 5: $E_4 = (b_{1,1}^k + \mu N_{PRB}^k; b_{1,2}^k; \dots; b_{1,m}^k; b_2^k; \dots; b_i^k + \mu'; \dots; b_m^k)$

$$(\mu, \mu') = \begin{cases} (1, -1) & \text{if a user } u_0 \text{ connects to network LTE in the cell } C_0 \text{ by} \\ & \text{disconnecting of the wireless Wi - Fi from the cell } C_i & \text{State}(E_{5,1}) \\ (-1, 1) & \text{if a user } u_0 \text{ disconnects of the network LTE from the cell } C_0 \\ & \text{by connecting to the wireless Wi - Fi in the cell } C_i & \text{State}(E_{5,2}) \\ (0, 0) & \text{Otherwise} & \text{State}(E_{5,3}) \end{cases}$$

- Stage 6: $E_5 = (b_{1,1}^k; b_{1,2}^k; \dots; b_{1,i}^k + \mu N_{PRB}^k; \dots; b_{1,m}^k; b_2^k; \dots; b_i^k + \mu'; \dots; b_m^k)$

$$(\mu, \mu') = \begin{cases} (1, -1) & \text{if a user } u_0 \text{ connects to the network LTE in the cell } C_i \text{ by} \\ & \text{disconnecting of the wireless Wi - Fi from the cell } C_i & \text{State}(E_{6,1}) \\ (-1, 1) & \text{if a user } u_0 \text{ disconnects of the network LTE from the cell } C_i \\ & \text{by connecting to the wireless Wi - Fi in the cell } C_i & \text{State}(E_{6,2}) \\ (0, 0) & \text{Otherwise} & \text{State}(E_{6,3}) \end{cases}$$

Transition rate

Proposition 1 Let us consider the heterogeneous system of networks in a stage $(E_p)_{1 \leq p \leq 6}$. The transition rate towards the state $E_{p,t}$ of the stage E_p is established by:

(1) In the zone C_0 of the cluster:

(2) In the cell C_i of the cluster:

$$\tau_{(E_0 \rightleftharpoons E_{1,1})} = \tau_{1,1} = (\lambda_{C_0}^{C(k)} + \lambda_{N_1}^{H(k)}) \cdot \left(\frac{b_{11}^k}{N_{PRB}^k} + 1 \right) \cdot \left(\frac{1}{\Delta_{C(k)}} + \eta_{C_1}^{\overline{C_1}} \right) \quad (8)$$



$$\tau_{(E_1 \Rightarrow E_{2,1})} = \tau_{2,1} = (\lambda_{C_i}^{C(k)} + \lambda_{C_i}^{C(k)} \cdot \mathbb{P}(N_1 \rightarrow N_i)) \cdot \left(\frac{b_{1i}^k}{N_{PRB}^k} + 1 \right) \cdot \frac{1}{\Delta_{C(k)}} \quad (9)$$

(3) From the zone C_0 to the cell C_i of cluster:

✓ By horizontal handover:

$$\tau_{(E_3 \Rightarrow E_{4,1})} = \tau_{4,1} = \left(\frac{b_{11}^k}{N_{PRB}^k} + 1 \right) \cdot \left(\frac{b_{1i}^k}{N_{PRB}^k} \right) \cdot \tau_{(C_0 \leftarrow C_i)}^{H(k)} \quad (10)$$

✓ By vertical handover:

$$\tau_{(E_4 \Rightarrow E_{5,1})} = \tau_{5,1} = (b_i^k) \cdot \left(\frac{b_{11}^k}{N_{PRB}^k} + 1 \right) \cdot \tau_{(C_0 \leftarrow C_i)}^{V(k)} \quad (11)$$

Proof

(1) When the number of units of busy bandwidth in one under zone C_1 of the network N_1 (LTE) crosses of b_{11} to $b_{11} + N_{PRB}^k$ then the variation rate of the units of busy bandwidth is defined by:

$$\frac{b_{11} + N_{PRB}^k}{N_{PRB}^k} = \frac{b_{11}}{N_{PRB}^k} + 1$$

If there is variation of units of busy bandwidth, it is because a user who is recently connected in it in sub-zone C_1 with a access demand rate $\lambda_{C_0}^{C(k)}$ or to make a horizontal handover $\lambda_{N_1}^{H(k)}$ and the residence time in this sub-zone C_1 is $\Delta_{C(k)}$ or the time of passage in this cell C_0 towards C_1 . The residence time is inversely proportional to the transition rate. The transition rate of the state E_0 to the state E_1 is equal to the product of these various rates calculated previously. So the found rate is:

$$(\lambda_{C_0}^{C(k)} + \lambda_{N_1}^{H(k)}) \cdot \left(\frac{b_{11}^k}{N_{PRB}^k} + 1 \right) \cdot \left(\frac{1}{\Delta_{C(k)}} + \eta_{C_0}^{C_1} \right)$$

(2) If a user u connects to the network N_1 (LTE) in a sub-zone C_i then the number of bandwidth units occupied in this sub-zone passes to b_{1i} $b_{1i} + N_{PRB}^k$. So bandwidth units rate variation occupied in this sub-zone is calculated by:

$$\frac{b_{1i} + N_{PRB}^k}{N_{PRB}^k} = \frac{b_{1i}}{N_{PRB}^k} + 1$$

This variation of the bandwidth units occupation rate is due to the fact that when a user is recently connected in this sub-zone C_i with an access request

rate $\lambda_{C_i}^{C(k)}$ rr disconnects from the network N_i by connecting to the network N_1 with a probability $\mathbb{P}(N_i \rightarrow N_1)$ applied to an access request rate $\lambda_{C_i}^{C(k)}$ in this sub-zone. The transition rate to the state $E_{2,1}$ of the stage E_2 being inversely proportional of stay time $\Delta_{C(k)}$ in this sub-zone C_i then it's found by making the product of the various rates calculated previously. So the rate $\tau_{(E_1 \Rightarrow E_{2,1})}$ found is equal to:

$$(\lambda_{C_i}^{C(k)} + \lambda_{C_i}^{C(k)} \cdot \mathbb{P}(N_i \rightarrow N_1)) \cdot \left(\frac{b_{1i}^k}{N_{PRB}^k} + 1 \right) \cdot \frac{1}{\Delta_{C(k)}}$$

(3) Besides, if a user u connects to the network N_1 (LTE) in sub-zone C_0 by disconnecting from the same network N_1 but in a zone C_i then the number of bandwidth units occupied in this sub-zone C_0 of the network N_1 (LTE) passes from b_{11} to $b_{11} + N_{PRB}^k$ and b_{1i} to $b_{1i} - N_{PRB}^k$ in sub-zone C_i . So the variation rate of the bandwidth units occupied in zones C_0 and C_i are respectively defined by:

$$\frac{b_{11} + N_{PRB}^k}{N_{PRB}^k} = \frac{b_{11}}{N_{PRB}^k} + 1$$

and

$$\frac{b_{1i}^k}{N_{PRB}^k}$$

As the user makes a horizontal handover (vertical resp.) of the zone C_i towards C_1 then the rate of transition is directly proportional at the horizontal handover rate (vertical resp.) $\tau_{(C_0 \leftarrow C_i)}^{H(k)}$ (resp. $\tau_{(C_0 \leftarrow C_i)}^{V(k)}$). So the transition rate of the state $E_{4,1}$ in the stage E_3 is equal to the product of these rates calculated previously. Thus the found rate is:

$$\left(\frac{b_{11}^k}{N_{PRB}^k} + 1 \right) \cdot \left(\frac{b_{1i}^k}{N_{PRB}^k} \right) \cdot \tau_{(C_0 \leftarrow C_i)}^{H(k)}$$

Respectively:

$$((b_i^k) \cdot \left(\frac{b_{11}^k}{N_{PRB}^k} + 1 \right) \cdot \tau_{(C_0 \leftarrow C_i)}^{V(k)})$$

□

The various transitions:

$$\checkmark \tau_{(E_0 \Rightarrow E_{1,1})} = (\lambda_{C_0}^{C(k)} + \lambda_{N_1}^{H(k)}) \cdot \left(\frac{b_{11}^k}{N_{PRB}^k} + 1 \right) \cdot$$

$$\left(\frac{1}{\Delta_{C(k)}} + \eta_{C_0}^{C_1} \right) \quad s.q / \sum_{j=1}^m (b_{1j} + N_{PRB}^k) \leq B_1^{uc}$$

$$\checkmark \tau_{(E_0 \Rightarrow E_{1,2})} = (\lambda_{C_0}^{C(k)} + \lambda_{N_1}^{H(k)}) \cdot \left(\frac{b_{11}^k}{N_{PRB}^k} \right) \cdot \left(\frac{1}{\Delta_{C(k)}} + \eta_{C_0}^{C_1} \right)$$

$$s.q / b_{11}^k \geq N_{PRB}^k$$



$$\begin{aligned}
\checkmark \quad \tau_{(E_1 \Rightarrow E_{2,1})} &= (\lambda_{C_i}^{C(k)} + \lambda_{C_i}^{C(k)} \cdot \mathbb{P}(N_1 > N_i)) \cdot \left(\frac{b_{1i}^k}{N_{\text{PRB}}^k} + 1 \right) \\
&\cdot \frac{1}{\Delta_{C(k)}} \cdot s.q / \left\{ \begin{array}{l} b_i = B_i \\ \sum_{j=1}^m (b_{1j} + N_{\text{PRB}}^k) \leq B_1^{\text{uc}} \end{array} \right. \\
\checkmark \quad \tau_{(E_1 \Rightarrow E_{2,2})} &= (\lambda_{C_i}^{C(k)} + \lambda_{C_i}^{C(k)} \cdot \mathbb{P}(N_1 > N_i)) \cdot \frac{b_{1i}^k}{N_{\text{PRB}}^k} \cdot \frac{1}{\Delta_{C(k)}} \\
&\cdot s.q / \left\{ \begin{array}{l} b_i = B_i \\ b_{1i}^k \geq N_{\text{PRB}}^k \end{array} \right. \\
\checkmark \quad \tau_{(E_2 \Rightarrow E_{3,1})} &= (\lambda_{C_i}^{C(k)} + \lambda_{C_i}^{C(k)} \cdot \mathbb{P}(N_1 > N_i)) \cdot (b_i^k + 1) \cdot \\
&\frac{1}{\Delta_{C(k)}} \cdot s.q / \left\{ \begin{array}{l} b_i \leq B_i \\ \sum_{j=1}^m (b_{1j} + N_{\text{PRB}}^k) = B_1^{\text{uc}} \end{array} \right. \\
\checkmark \quad \tau_{(E_2 \Rightarrow E_{3,2})} &= (\lambda_{C_i}^{C(k)} + \lambda_{C_i}^{C(k)} \cdot \mathbb{P}(N_1 > N_i)) \cdot (b_i^k) \cdot \frac{1}{\Delta_{C(k)}} \\
&\cdot s.q / \left\{ \begin{array}{l} b_i \geq 1 \\ \sum_{j=1}^m (b_{1j} + N_{\text{PRB}}^k) = B_1^{\text{uc}} \end{array} \right. \\
\checkmark \quad \tau_{(E_3 \Rightarrow E_{4,1})} &= \left(\frac{b_{11}^k}{N_{\text{PRB}}^k} + 1 \right) \cdot \left(\frac{b_{1i}^k}{N_{\text{PRB}}^k} \right) \cdot \tau_{(C_i > C_0)}^{H(k)} \\
&\cdot s.q / \left\{ \begin{array}{l} b_i = B_i \\ \sum_{j=1}^m (b_{1j} + N_{\text{PRB}}^k) \leq B_1^{\text{uc}} \\ b_{1i}^k \geq N_{\text{PRB}}^k \end{array} \right. \\
\checkmark \quad \tau_{(E_3 \Rightarrow E_{4,2})} &= \left(\frac{b_{11}^k}{N_{\text{PRB}}^k} \right) \cdot \left(\frac{b_{1i}^k}{N_{\text{PRB}}^k} + 1 \right) \cdot \tau_{(C_i > C_0)}^{H(k)} \\
&\cdot s.q / \left\{ \begin{array}{l} b_i = B_i \\ \sum_{j=1}^m (b_{1j} + N_{\text{PRB}}^k) \leq B_1^{\text{uc}} \\ b_{1i}^k \geq N_{\text{PRB}}^k \end{array} \right. \\
\checkmark \quad \tau_{(E_4 \Rightarrow E_{5,1})} &= (b_i^k) \cdot \left(\frac{b_{11}^k}{N_{\text{PRB}}^k} + 1 \right) \cdot \tau_{(C_i > C_0)}^{V(k)} \\
&\cdot s.q / \left\{ \begin{array}{l} b_i \geq 1 \\ \sum_{j=1}^m (b_{1j} + N_{\text{PRB}}^k) \leq B_1^{\text{uc}} \end{array} \right. \\
\checkmark \quad \tau_{(E_4 \Rightarrow E_{5,2})} &= (b_i^k + 1) \cdot \left(\frac{b_{11}^k}{N_{\text{PRB}}^k} \right) \cdot \tau_{(C_i > C_0)}^{V(k)} \\
&\cdot s.q / \left\{ \begin{array}{l} b_i \leq B_i \\ b_{1i}^k \geq N_{\text{PRB}}^k \end{array} \right. \\
\checkmark \quad \tau_{(E_5 \Rightarrow E_{6,1})} &= \left(\frac{b_{1i}^k}{N_{\text{PRB}}^k} + 1 \right) \cdot (b_i^k) \\
&\cdot s.q / \left\{ \begin{array}{l} b_i \geq 1 \\ \sum_{j=1}^m (b_{1j} + N_{\text{PRB}}^k) \leq B_1^{\text{uc}} \end{array} \right. \\
\checkmark \quad \tau_{(E_5 \Rightarrow E_{6,2})} &= \left(\frac{b_{1i}^k}{N_{\text{PRB}}^k} \right) \cdot (b_i^k + 1) \cdot s.q / \left\{ \begin{array}{l} b_{1i}^k \geq N_{\text{PRB}}^k \\ b_i \leq B_i \end{array} \right.
\end{aligned}$$

Evaluations of optimal performances

We estimate the selection strategy performances of a network based on the bit rate related to the parameters such as the blocking probability and the connections quality.

Average bit rate of the system in a sub-zone C_i

The bit rate received in a zone Z_i from the network N_1 is in function of the number of present units of bandwidth in networks N_1 and N_i . By denoting $D_1^{\text{avg}}(E)$ the average bit

rate value received from the network and $\mathbb{P}(E)$ the probability of system balance state so the total average bit rate value received in a zone Z_i from the network N_1 is given by:

$$D_{N_1}^{\text{tot}} = \sum_{k=1}^s (\lambda_{C_i}^{C(k)} + \mathbb{P}(C_i > C_1)) \cdot \mathbb{P}(E) \cdot D_1^{\text{avg}}(E) \quad (12)$$

$$s.q / \sum_{j=1}^m (b_{1j} + N_{\text{PRB}}^k) \leq B_1^{\text{uc}}$$

with:

$$D_1^{\text{avg}}(E) = D_1^{\text{avg}} \cdot \left(1 - \Lambda \cdot \sqrt{\frac{\sum_{j=1}^m (b_{1j} + b_j)}{B_1^{\text{uc}} + B_i}} \right)$$

where the average instantaneous bit rate $D_1^{\text{avg}}(E)$ is defined by the product of a sub-carrier bandwidth and modulation(numbers of modulated sub-carriers):

$$D_1^{\text{avg}} = B_{\text{sp}} \times N_{\text{sub}} \quad (13)$$

where $N_{\text{sub}} = K \times B \times E_i \times (1 - \text{BLER}_i)$ such as:

- K is the frequencies number;
- B is the numbers of symbols per second;
- E_i is the modulation efficiency;
- BLER_i is the Block error rate in a sub-cell C_i ;

As result we have:

$$\bar{D}_{C_i} = \frac{D_{N_1}^{\text{tot}}}{\eta_{C_i}^{c(k)}} \quad (14)$$

Blocking probability et connections losses in sub-cell

C_i

We have calculated the mean blocking probability of connections in a sub-zone C_i in function of an equilibrium state probability $\mathbb{P}(E)$ of the system. Indeed, we added the system states probability where the numbers of units of busy bandwidth is higher than these available in the network N_1 by formula:

$$\mathbb{P}_{N_1}^B = \sum_{k=1}^s (\lambda_{C_i}^{C(k)} + \mathbb{P}(C_i > C_1)) \cdot \mathbb{P}(E) \cdot \mathbb{P}_1^B(E) \quad (15)$$

$$s.q / \sum_{j=1}^m (b_{1j} + N_{\text{PRB}}^k) > B_1^{\text{uc}}$$

with

$$\mathbb{P}_1^B(E) = \mathbb{P}_1^B \cdot \left(1 - \Theta \cdot \sqrt{\frac{\sum_{j=1}^m (b_{1j} + b_j)}{B_1^{\text{uc}} + B_i}} \right)$$



Table 4 LTE network parameters test

Test	Value
Modulation	16 QAM
Symbols	6
Efficiency	1.4766
Sub-carriers number	72
Bandwidth	1.4 MHz

Table 5 System network parameters test

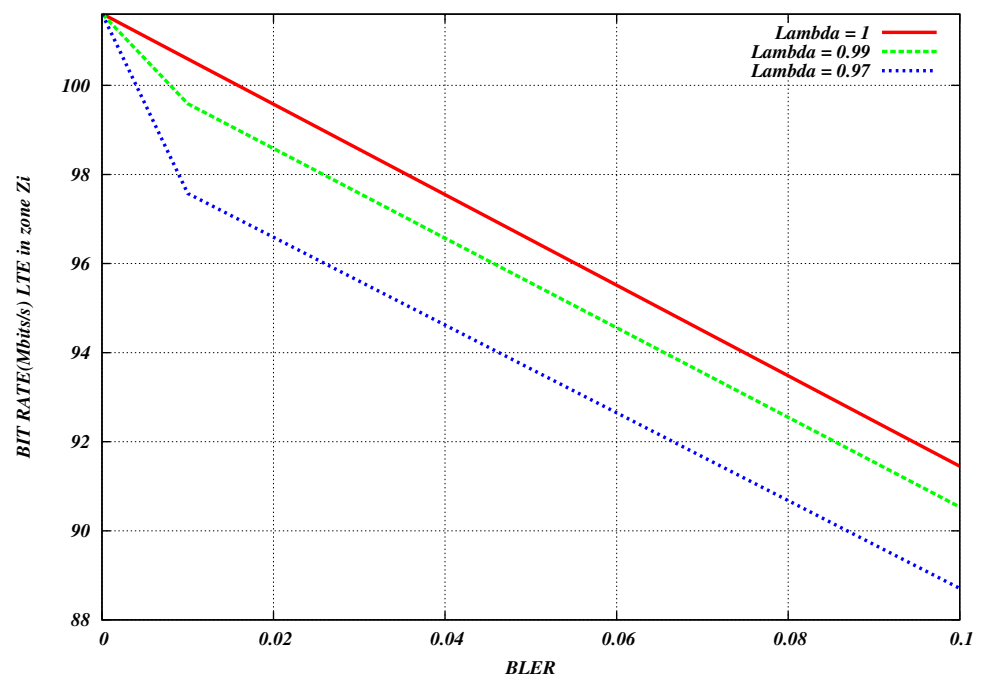
Parameters test	Value
Average access rate to a service	70 %
Average handover rate	60 %
Average state balance rate	80 %
Bandwidth units in LTE	60
Noise power	−174 dBm/Hz
Signal power	400 dBm
Services number	Two services unicast
Cell radius	600 m Z_1 , 200 m Z_2
Distance between area centers	300 m

$$\mathbb{P}_1^B = \frac{\frac{\rho_{C_i}^s}{s!}}{\sum_{k=1}^s \frac{\rho_{C_i}^k}{k!}} \quad (16)$$

where s is the number of available services in the cluster and

$$\rho_{C_i} = \sum_{k=1}^s \left(\frac{\lambda_{C_i}^{c(k)}}{\lambda_{C_i}^c} \right)$$

is the probability that a user is blocked in the sub-zone C_i .

Fig. 3 Average bit rate of network N_1 in the cell C_2 in function of the BLER

Numerical tests

To test the theoretical results that we obtained, we simplified our field of study in a service cell C_1 covered by the network N_1 (LTE) in which we implanted a wireless N_2 (Wi-Fi) in sub-cell C_2 of C_1 . By means of the simulator NS3 [9] and working with the parameters below we managed to obtain satisfactory results as show by the obtained curves.

We have supposed that the data used for the test are the ones relative to the characteristics of the network LTE. So, we chose the modulation 16QAM of the network LTE the number of symbols which is six (06), Efficiency is equal 1.4766, sub-carriers number is 72 and as the bandwidth of the network LTE varies between 1.4 and 20 MHz then we have worked with the minimal value (1.4MHz) as indicated in Table 4; refer to Jabban [8].

However, the parameters chosen for the heterogeneous mobile and wireless system of networks are supposed to allow us to have satisfactory results for the numerical results. These parameters are given in Table 5.

The main parameters of both available networks are illustrated in the Table 5. We suppose that the coverage radius of networks N_1 and N_2 are respectively equal to 600 M and 200 m. We also suppose that the power transmitted by N_1 is equal to 400 dBm. The noise power is supposed equal to −174 dBm/Hm for the network LTE. We also suppose that the average access requests rates for a service, handover and the system balance state are, respectively, 70, 60 and 80 %. The number of units of bandwidth is supposed equal to 60. Finally the number of unicasts services is fixed to 2.



Fig. 4 Average bit rate of the network N_1 in the zone C_2 in function of the busy bandwidth rate

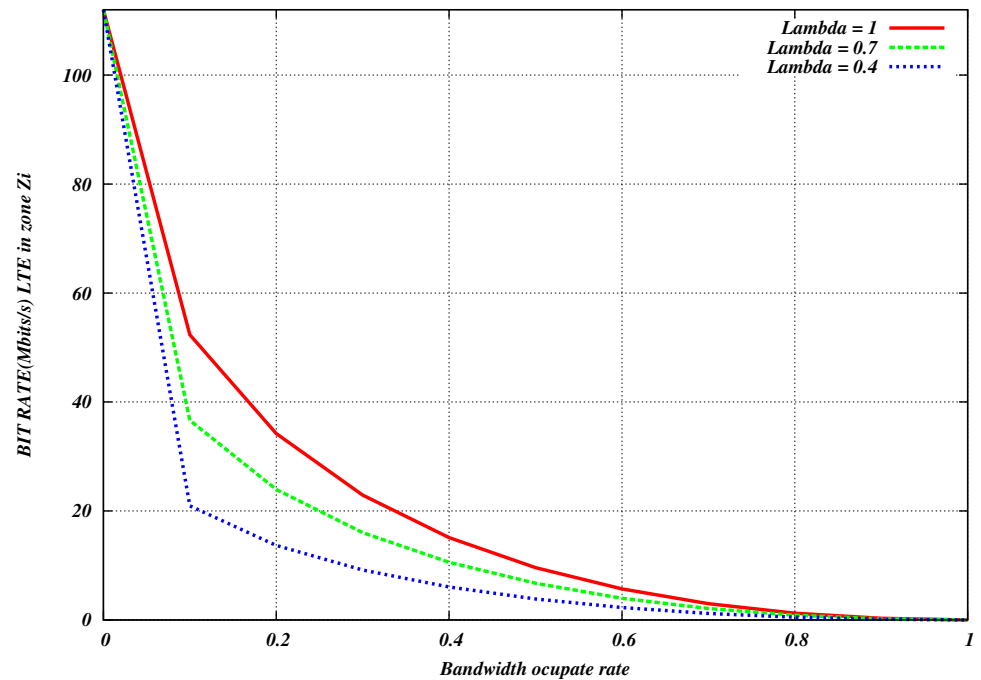
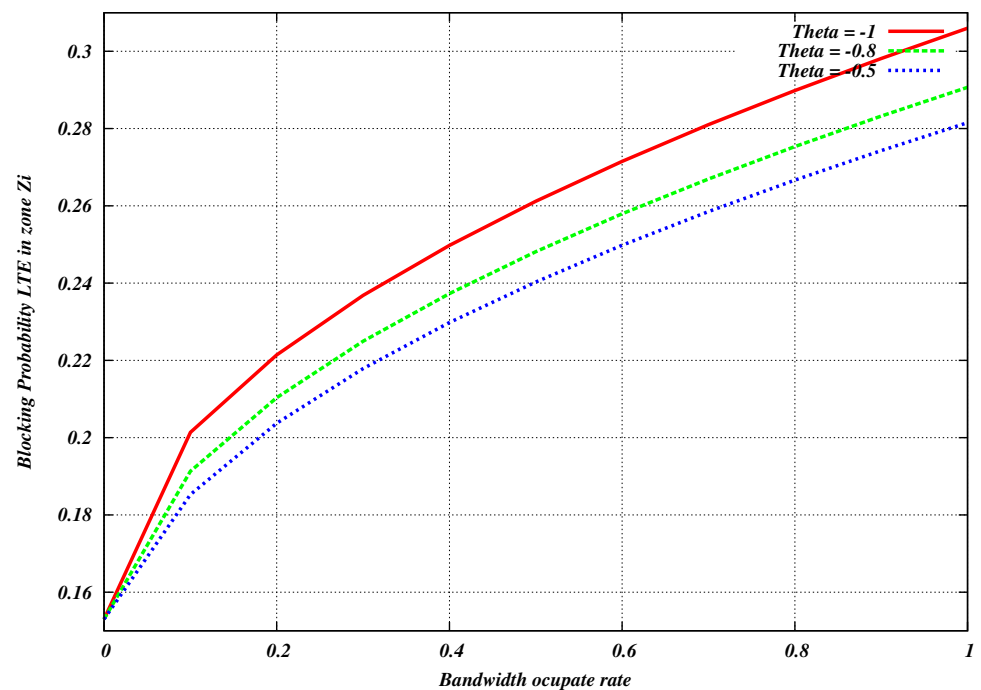


Fig. 5 Blocking probability in the zone C_2 in function of the busy bandwidth rate

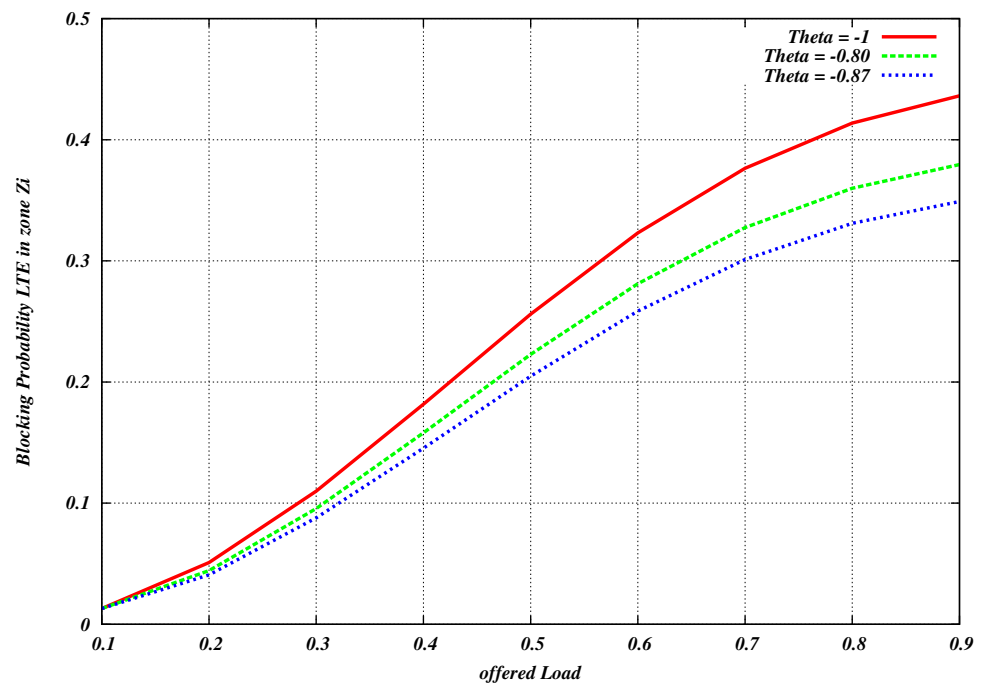


We have calculated the average bit rate value $D_{N_1}^{\text{tot}}$ received by a user u_0 from the network N_1 in the cell C_2 . The selection strategy is based on the most high bit rate, in a sub-cell C_2 , the user chooses the unicasts services of the networks N_1 or N_2 according to the selection technique established. Through the obtained results, we have noticed that a small modification of the bit rate occurs

due to the network saturation. This modification is represented by a parameter Λ which indicates a very high sensitivity of the bit rate because of the network congestion. For a low modification of the sensitivity factor Λ for example from 1 to 0.99 we obtained a net fall of the bit rate as illustrate by Fig. 2 in function of the BLER in the sub-cell C_2 . This factor also informs about the



Fig. 6 Blocking probability in the zone C_2 in function of the offered load rate



maximal decrease of the bit rate because of the network congestion.

However the change of the factor Λ acts less on the bit rate when it is estimated in function of the occupation bandwidth rate by fixing the BLER as indicated by Fig. 3; BLER is fixed to 50 %.

Besides, we have estimated the network performances related to the average blocking probability and connections losses to the services in the sub-cell C_2 . The obtained results depend on a sensitivity factor Θ as represented by Figs. 4 and 5. For sensitivity parameters $\Theta \in \{1; 0.8; 0.5\}$ we analyzed the blocking probability in the sub-cell C_2 in function of the busy bandwidth rate (Fig. 4) then to the offered load by the traffic (Fig. 5). Indeed, seen the satisfactory obtained results, we realize that the blocking probability do not pass 40 % when they are determined with the occupation bandwidth rate. They do not reach either the level 50 % when they are estimated in function of the offered load by the traffic whatever is the given sensitivity factor (Fig. 6).

Conclusion

At the end of our analysis on the system integration performances of new generation wireless and mobile networks, we found a factor which remains very sensitive to the variations of the bit rate received in a sub-cell C_i when it is calculated in function of the BLER in this sub-zone.

Besides, this sensitivity factor remains so determining for the blocking probability theory in a sub-zone C_i by means of the busy bandwidth rate or the offered load traffic rate. The satisfactory results obtained on the system performances of wireless and mobile networks such the LTE and the Wi-Fi based on the bit rate allowed us to discover the dynamic fluctuations system. The parameters related to the bit rate such as the blocking probability are estimated with lower rates the bar 40 %.

In our future works, we intend to calculate the same sensitivity factor when we consider the performances of the system related to the SINR. We planned also to take into account the number of users having consumed these numbers of busy bandwidth units. This will allow us to encircle better the congestion and dis-congestion rates of the heterogeneous networks.

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